

Krylov subspace methods for linear systems with tensor product structure

Christine Tobler

Seminar for Applied Mathematics,
ETH Zürich

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Outline

- 1 Introduction
- 2 Basic Algorithm
- 3 Convergence bounds
- 4 Solving the compressed equation
- 5 Numerical experiments

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Linear system with tensor product structure

We consider the linear system

$$\mathcal{A}x = b$$

with

$$\mathcal{A} = \sum_{s=1}^d I_{n_1} \otimes \cdots \otimes I_{n_{s-1}} \otimes A_s \otimes I_{n_{s+1}} \otimes \cdots \otimes I_{n_d},$$

$$b = b_1 \otimes \cdots \otimes b_d,$$

$$A_s \in \mathbb{R}^{n_s \times n_s} \text{ positive definite, } b_s \in \mathbb{R}^{n_s}.$$

Example for 3 dimensions:

$$(A_1 \otimes I \otimes I + I \otimes A_2 \otimes I + I \otimes I \otimes A_3)x = b_1 \otimes b_2 \otimes b_3$$

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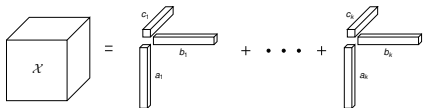
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Tensor decompositions

CP decomposition:

$$\text{vec}(\mathcal{X}) = \sum_{r=1}^k \mathbf{a}_r \otimes \mathbf{b}_r \otimes \mathbf{c}_r$$

$$\mathbf{a}_r \in \mathbb{R}^m, \mathbf{b}_r \in \mathbb{R}^n, \mathbf{c}_r \in \mathbb{R}^p$$



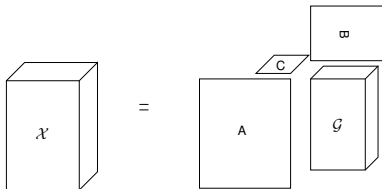
Tucker decomposition:

$$\text{vec}(\mathcal{X}) = \sum_{i=1}^{\tilde{m}} \sum_{s=1}^{\tilde{n}} \sum_{l=1}^{\tilde{p}} \mathcal{G}_{ijl} \mathbf{a}_i \otimes \mathbf{b}_j \otimes \mathbf{c}_l$$

$$= (\mathbf{A} \otimes \mathbf{B} \otimes \mathbf{C}) \text{vec}(\mathcal{G})$$

$$\mathcal{G} \in \mathbb{R}^{\tilde{m} \times \tilde{n} \times \tilde{p}},$$

$$\mathbf{a}_i \in \mathbb{R}^m, \mathbf{b}_j \in \mathbb{R}^n, \mathbf{c}_l \in \mathbb{R}^p$$



About this system

The eigenvalues of the matrix \mathcal{A} are given by all possible sums

$$\lambda_{i_1}^{(1)} + \lambda_{i_2}^{(2)} + \dots + \lambda_{i_d}^{(d)}$$

where $\lambda_{i_s}^{(s)}$ denotes an eigenvalue of A_s . For A_s positive definite, the system has a unique solution.

Note that x and b are vector representations of tensors in $\mathbb{R}^{n_1 \times \dots \times n_d}$, and b represents a rank-one tensor.

A tensor arising from the discretization of a sufficiently smooth function f can be approximated by a short sum of rank-one tensors. By superposition, we can insert such a tensor as right-hand side into our system.

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Tensorized Krylov subspaces

A Krylov subspace is defined as

$$\mathcal{K}_k(\mathbf{A}, \mathbf{b}) = \text{span}\{\mathbf{b}, \mathbf{A}\mathbf{b}, \dots, \mathbf{A}^{k-1}\mathbf{b}\}.$$

We define a tensorized Krylov subspace as

$$\mathcal{K}_{\mathfrak{K}}^{\otimes}(\mathcal{A}, \mathbf{b}) := \text{span}(\mathcal{K}_{k_1}(\mathbf{A}_1, \mathbf{b}_1) \otimes \dots \otimes \mathcal{K}_{k_d}(\mathbf{A}_d, \mathbf{b}_d)).$$

Note that

$$\mathcal{K}_{k_0}(\mathcal{A}, \mathbf{b}) \subset \mathcal{K}_{\mathfrak{K}}^{\otimes}(\mathcal{A}, \mathbf{b})$$

for $\mathfrak{K} = (k_0, \dots, k_0)$.

⇒ Tensorized Krylov subspaces are richer than standard Krylov subspaces.

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Reminder: CG method ($Ax = b$)

The best approximation of x in $\mathcal{K}_k(A, b)$ is defined by:

$$\|x_k - x\|_A = \min_{\tilde{x} \in \mathcal{K}_k(A, b)} \|\tilde{x} - x\|_A.$$

Find U_k with (orthonormal) columns that span the Krylov subspace $\mathcal{K}_k(A, b)$. Set $x_k = U_k y$, then y is the solution of the compressed system

$$H_k y = \tilde{b}, \quad H_k = U_k^\top A U_k, \quad \tilde{b} = U_k^\top b.$$

Convergence bound ($\kappa = \kappa_2(A)$)

$$\|x_k - x\|_A \leq C(A, b, \kappa) \left(\frac{\sqrt{\kappa} - 1}{\sqrt{\kappa} + 1} \right)^k.$$

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Tensorized Krylov: $\mathcal{A} \operatorname{vec}(\mathcal{X}) = b$

Applying the Arnoldi method to A_s, b_s results in U_s, H_s with

$$U_s^\top A_s U_s = H_s,$$

U_s column-orthogonal, H_s upper Hessenberg matrix.

The columns of U_s span $\mathcal{K}_{k_s}(A_s, b_s)$, and similarly the columns of $U := U_1 \otimes \cdots \otimes U_d$ span $\mathcal{K}_{\mathbb{R}}^{\otimes}(\mathcal{A}, b)$.

Solve the compressed system

$$\mathcal{H}y = \tilde{b}$$

with $x_{\mathbb{R}} = Uy$, $\tilde{b} = U^\top b$ and $\mathcal{H} = U^\top \mathcal{A}U$.

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Solve the compressed system

$$\mathcal{H}y = \tilde{b}$$

with $x_{\mathbb{R}} = \mathcal{U}y$, $\tilde{b} = \mathcal{U}^\top b$ and $\mathcal{H} = \mathcal{U}^\top \mathcal{A} \mathcal{U}$.

Compressed system

The structure of $\mathcal{H}y = \tilde{b}$ corresponds to that of $\mathcal{A}x = b$:

$$\mathcal{H} = \sum_{s=1}^d I_{k_1} \otimes \cdots \otimes I_{k_{s-1}} \otimes H_s \otimes I_{k_{s+1}} \otimes \cdots \otimes I_{k_d}$$
$$\tilde{b} = \tilde{b}_1 \otimes \cdots \otimes \tilde{b}_d = \|b\|_2 (\mathbf{e}_1 \otimes \cdots \otimes \mathbf{e}_1)$$

For dense core tensor y , $x_{\mathcal{R}}$ is a tensor in Tucker decomposition:

$$x_{\mathcal{R}} = (U_1 \otimes \cdots \otimes U_d)y$$

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Convergence, s.p.d. case (1)

Similarly to CG, we have:

$$\|x_{\mathcal{K}} - x\|_{\mathcal{A}} = \min_{\tilde{x} \in \mathcal{K}_{\mathcal{K}}^{\otimes}(\mathcal{A}, b)} \|\tilde{x} - x\|_{\mathcal{A}}$$

Every vector in a Krylov subspace $\mathcal{K}_{k_s}(A_s, b_s)$ can be seen as $p(A_s)b_s$, with p a polynomial of order at most $k_s - 1$. A similar thought reduces the bound calculation to the min-max problem

$$E_{\Omega}(\mathcal{K}) := \min_{p \in \Pi_{\mathcal{K}}^{\otimes}} \left\| p(\lambda_1, \dots, \lambda_d) - \frac{1}{\lambda_1 + \dots + \lambda_d} \right\|_{\Omega},$$

where $\Pi_{\mathcal{K}}^{\otimes}$ is a space of multivariate polynomials, and Ω contains the eigenvalues $(\lambda_1, \dots, \lambda_d)$, with $\lambda_i \in [\lambda_{\min}(A_i), \lambda_{\max}(A_i)]$.

Convergence, s.p.d. case (2)

Inserting an upper bound for $E_{\Omega}(\hat{\mathcal{K}})$, we find

$$\|x_{\hat{\mathcal{K}}} - x\|_{\mathcal{A}} \leq \sum_{s=1}^d C(\mathcal{A}, b, \kappa_S) \left(\frac{\sqrt{\kappa_S} - 1}{\sqrt{\kappa_S} + 1} \right)^{k_S},$$

with $\kappa_S = 1 + \frac{\lambda_{\max}(A_S) - \lambda_{\min}(A_S)}{\lambda_{\min}(\mathcal{A})}$.

For the case A_S, k_S constant:

$$\|x_{\hat{\mathcal{K}}} - x\|_{\mathcal{A}} \leq C(\mathcal{A}, b, d) \left(\frac{\sqrt{\kappa} - 1}{\sqrt{\kappa} + 1} \right)^k,$$

with $\kappa = \frac{d-1}{d} + \frac{\kappa_2(\mathcal{A})}{d}$.

Note that the convergence rate improves with increasing dimension.

Convergence, non-symmetric positive definite case

General convergence bound:

$$\|x_{\hat{\mathcal{R}}} - x\|_2 \leq \sum_{s=1}^d \int_0^{\infty} e^{-\hat{\alpha}_s t} \|U_s e^{-tH_s} e_1 - e^{-tA_s} b_s\|_2 dt$$

with $\hat{\alpha}_s := \sum_{j \neq s} \alpha_j$ and $\alpha_j = \lambda_{\min}(A_j + A_j^T)/2$.

The bound on $\|U_s e^{-tH_s} e_1 - e^{-tA_s} b_s\|_2$ will depend on additional knowledge on A_s .

For example, when the field of values of each matrix A_s is contained in a known ellipse, an explicit convergence bound can be found.

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Grasedyck's method, system $\mathcal{H}y = \tilde{b}$

For \mathcal{H} positive definite:

$$\mathcal{H}^{-1} = \int_0^{\infty} \exp(-t\mathcal{H}) dt$$

The exponential of \mathcal{H} has a Kronecker product structure, too:

$$\exp(-t\mathcal{H}) = \exp\left(-t \sum_{s=1}^d \hat{H}_s\right) = \prod_{s=1}^d \exp(-t\hat{H}_s) = \bigotimes_{s=1}^d \exp(-tH_s),$$

Approximation y_m :

$$y_m = \sum_{j=1}^m \omega_j \bigotimes_{s=1}^d \exp\left(-\alpha_j H_s\right) \tilde{b}_s,$$

with certain coefficients α_j, ω_j .

Coefficients of the exponential sum (1)

The coefficients α_j, ω_j should minimize

$$\sup_{z \in \Lambda(\mathcal{H})} \left| \frac{1}{z} - \sum_{j=1}^m \omega_j e^{-\alpha_j z} \right|$$

Case 1: \mathcal{H} symmetric and condition number known

The eigenvalues of \mathcal{H} are real: $\Lambda(\mathcal{H}) \subset [\lambda_{\min}, \lambda_{\max}]$. There are coefficients α_j, ω_j s.t.

$$\|y - y_m\|_2 \leq C(\mathcal{H}, \tilde{b}) \exp\left(\frac{-m\pi^2}{\log(8\kappa_2(\mathcal{H}))}\right).$$

These coefficients may be found using a variant of the Remez algorithm. Tabellated values have been made available by Hackbusch.

Coefficients of the exponential sum (2)

Case 2: Nonsymmetric \mathcal{H} and/or unknown condition number

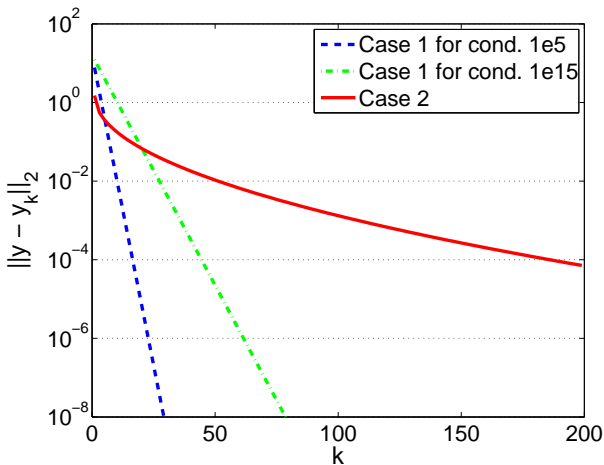
There is an explicit formula for α_j, ω_j , where the eigenvalues of \mathcal{H} only need to have positive real part.

$$\|y - y_{2m+1}\|_2 \leq C(\mathcal{H}, \tilde{\mathbf{b}}) \exp(\mu/\pi) \exp(-\sqrt{m}),$$

where $\mu = \max\{|\Im m(\Lambda(\mathcal{H}))|\}$.

The convergence is significantly slower than for Case 1.

Theoretical convergence bounds for the coefficients



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Symmetric example: Poisson Equation

Finite difference discretization of the Poisson equation in d dimensions:

$$\begin{aligned}\Delta u &= f \quad \text{in } \Omega = [0, 1]^d \\ u &= 0 \quad \text{on } \Gamma := \partial\Omega,\end{aligned}$$

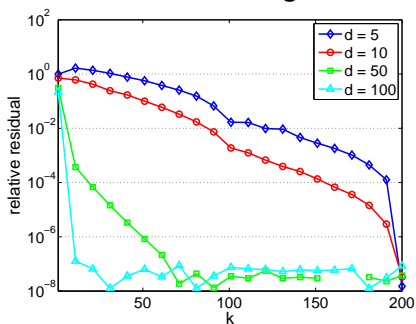
where the right-hand side f is a separable function.

$$A_s = \frac{1}{h^2} \begin{pmatrix} 2 & -1 & & & \\ -1 & 2 & \ddots & & \\ & \ddots & \ddots & \ddots & \\ & & & -1 & 2 \\ & & & & -1 & 2 \end{pmatrix}, b_s : \text{random numbers.}$$

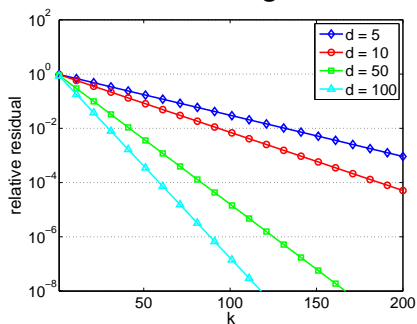
The approximation error is measured by *relative residual*, $\frac{\|Ax_{\mathbb{R}} - b\|_2}{\|b\|_2}$.

Convergence for system size 200^d

Numerical convergence

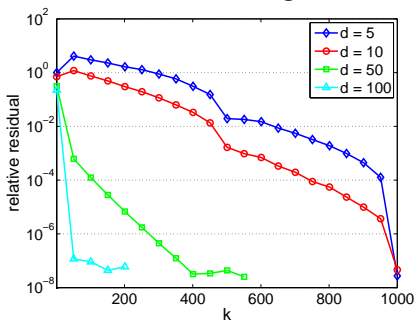


Theoretical convergence rate

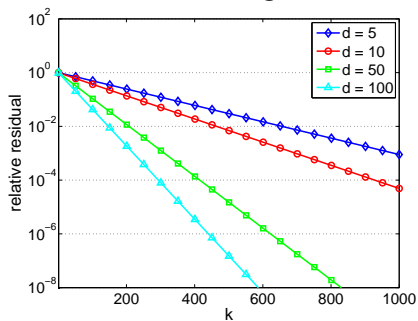


Convergence for system size 1000^d

Numerical convergence



Theoretical convergence rate



Extended Krylov subspaces

Using *extended Krylov subspaces*

$$\tilde{\mathcal{K}}_{k_s}(A_S, b_S) := \text{span}(\mathcal{K}_{k_s}(A_S, b_S) \cup \mathcal{K}_{k_s+1}(A_S^{-1}, b_S)),$$

the algorithm works analogously.

Convergence bound (A_S, k_S constant):

$$\|x_{\tilde{\mathcal{R}}} - x\|_2 \leq C(\mathcal{A}, b, d) \left(\frac{\sqrt{\tilde{\kappa}} - 1}{\sqrt{\tilde{\kappa}} + 1} \right)^k.$$

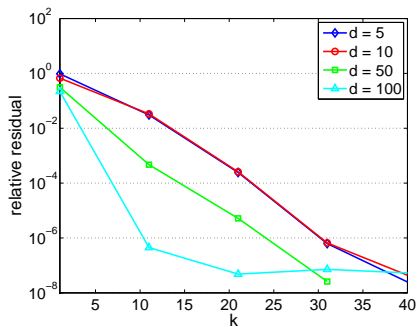
The convergence rate depends on $\tilde{\kappa}$:

$$\tilde{\kappa} \approx \frac{d-1}{d} + \frac{1}{d^2} (d-1)^{\frac{d-1}{d}} \kappa_2(\mathcal{A})^{\frac{d-1}{d}}$$

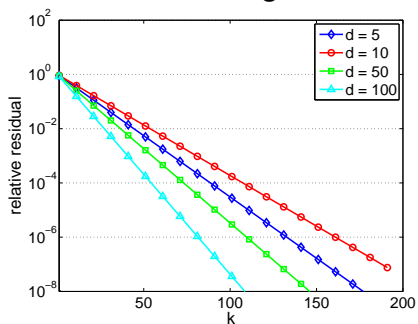
$$d = 2 : \tilde{\kappa} = \frac{1}{2} + \frac{\sqrt{\kappa_2(\mathcal{A})}}{4}, \quad d \gg 0 : \tilde{\kappa} \approx \frac{d-1}{d} + \frac{\kappa_2(\mathcal{A})}{d}.$$

Extended Krylov, system size 200^d

Numerical convergence

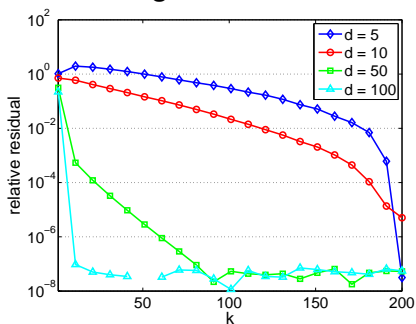


Theoretical convergence rate

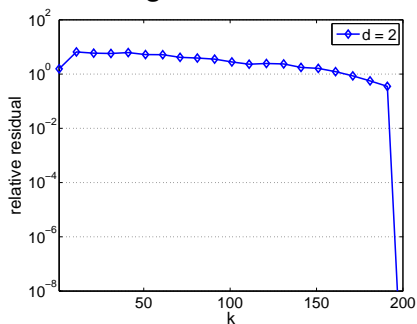


Non-symmetric case, system size 200^d

Convergence for $c = 10$



Convergence for $c = 100$



Conclusions

- An efficient algorithm to calculate a low-rank approximation of the solution tensor
- The computational complexity is linear in the number of dimensions
- Only matrix-vector operations with the full system matrices are required
- A theoretical convergence bound was found.

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Thank you for your attention!

Literature

D. Kressner, C. Tobler, “Krylov subspace methods for linear systems with tensor product structure”, *Research report No. 2009-16, SAM, ETH Zurich*, 2009.

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